

# Shuffle Theorems and Sandpiles

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## The Shuffle Theorem

The *shuffle theorem* of Carlsson and Mellit gives a combinatorial expression for coeffs of the symmetric function  $\nabla e_n$  in terms of a  $q, t$ -weighted count of labeled lattice paths:

**Theorem (Carlsson–Mellit [1] + Loehr–Remmel [2]):**

For all  $\mu, \nu$  compositions with  $|\mu| + |\nu| = n$ ,

$$\langle \nabla e_n, e_\mu h_\nu \rangle = \sum_{D \in \text{PF}(\mu; \nu)} q^{\text{area}(D)} t^{\text{pmaj}(D)},$$

where  $\text{PF}(\mu; \nu)$  is a certain class of parking functions (viewed as lattice paths).

$\text{area}(D) = \#$  whole squares between  $D$  and main diagonal.

$\text{pmaj}(D)$  is significantly more complicated.

## Graphs

For all  $\mu, \nu$  with  $|\mu| + |\nu| = n$ ,  $G_{\mu, \nu}$  is a graph on vertex set  $[n]$  with:

- A clique  $K_{\mu_i}$  for each component  $\mu_1, \dots, \mu_{\ell(\mu)}$  of  $\mu$ .
- An independent set  $\bar{K}_{\nu_i}$  for each component  $\nu_1, \dots, \nu_{\ell(\nu)}$  of  $\nu$ .
- All possible edges between these blocks.

We fix a particular labeling of the vertices:

- The vertices of  $\mu_i$  are  $n - (\mu_1 + \dots + \mu_{i-1})$  through  $n - (\mu_1 + \dots + \mu_i) + 1$ .
- The vertices of  $\nu_i$  are  $n - |\mu| - (\nu_1 + \dots + \nu_{i-1})$  through  $n - |\mu| - (\nu_1 + \dots + \nu_i) + 1$ .

Attach a *sink vertex* 0 to each vertex of  $G_{\mu, \nu}$  to get  $\hat{G}_{\mu, \nu}$ .

Running example:  $\hat{G}_{(4,3),(2,3)}$

## Sandpiles

*Sandpile configuration*  $\kappa$  on  $G$ : assign a nonneg. number of “grains of sand” to each nonsink vertex of  $G$  (sink has unlimited sand).

If  $\kappa(v) \geq \deg(v)$ ,  $v$  can *topple* and give one grain to each of its neighbors.  $\phi_v(\kappa) =$  resulting config. If  $v$  can't topple,  $\kappa$  is *stable* at  $v$ .

Stable config  $\kappa$  is *recurrent* if, starting from  $\phi_{\text{sink}}(\kappa)$ , nonsink verts can be ordered so that if we topple in this order the config always stays nonnegative.

## Sorted Recurrent Configurations

Sandpiles on  $\hat{G}_{\mu, \nu}$  admit an action of the Young subgroup  $\mathfrak{S}_\mu \times \mathfrak{S}_\nu$ .

Going mod this action  $\iff$  sorting the values of  $\kappa$  on each block of vertices corresponding to entries of  $\mu, \nu$ .

Recurrent config  $\kappa$  is *sorted recurrent* if:

- $\kappa$  is weakly decreasing inside each clique  $K_{\mu_i}$ .
- $\kappa$  is weakly increasing inside each independent block  $\bar{K}_{\nu_i}$ .

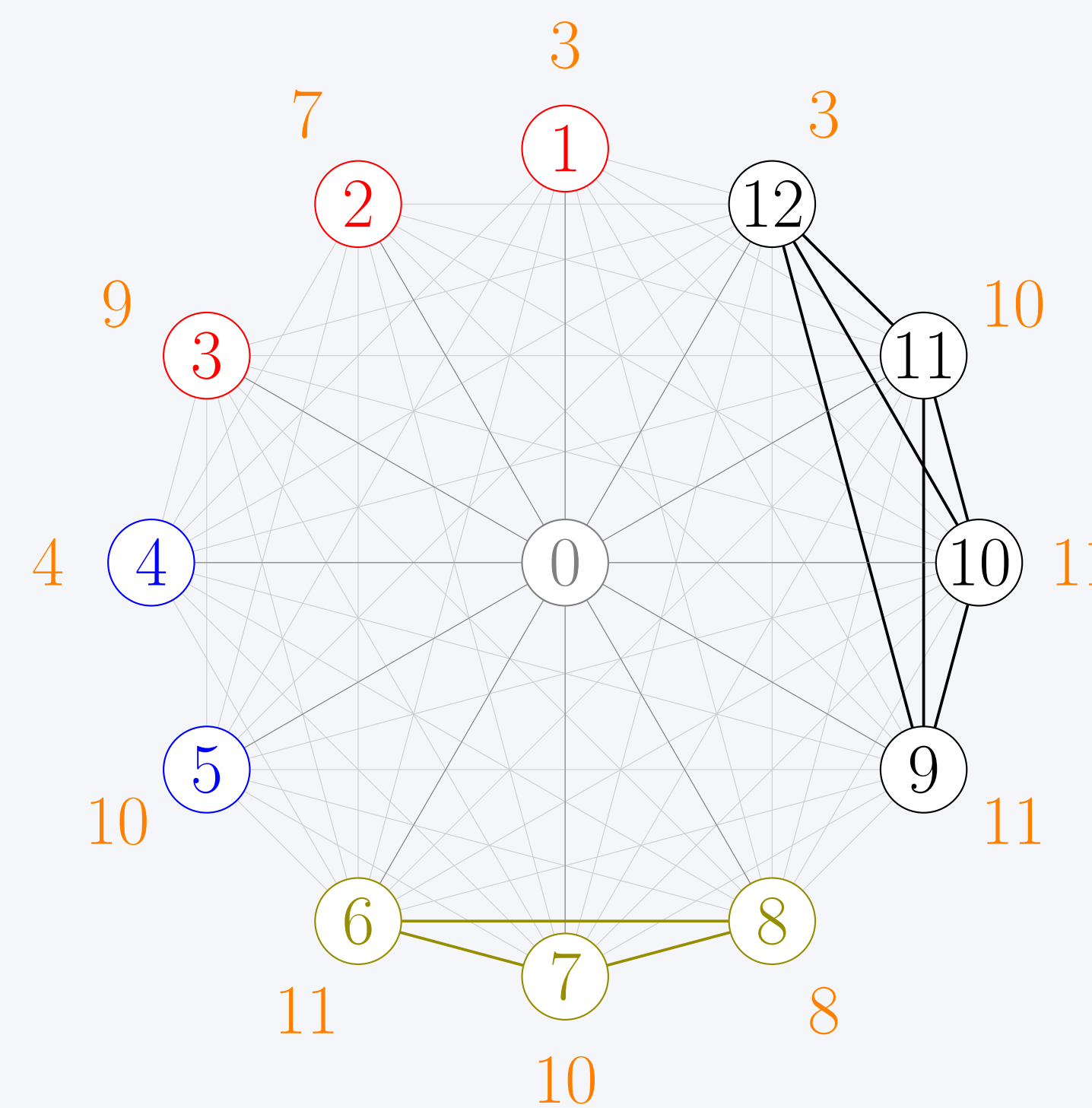
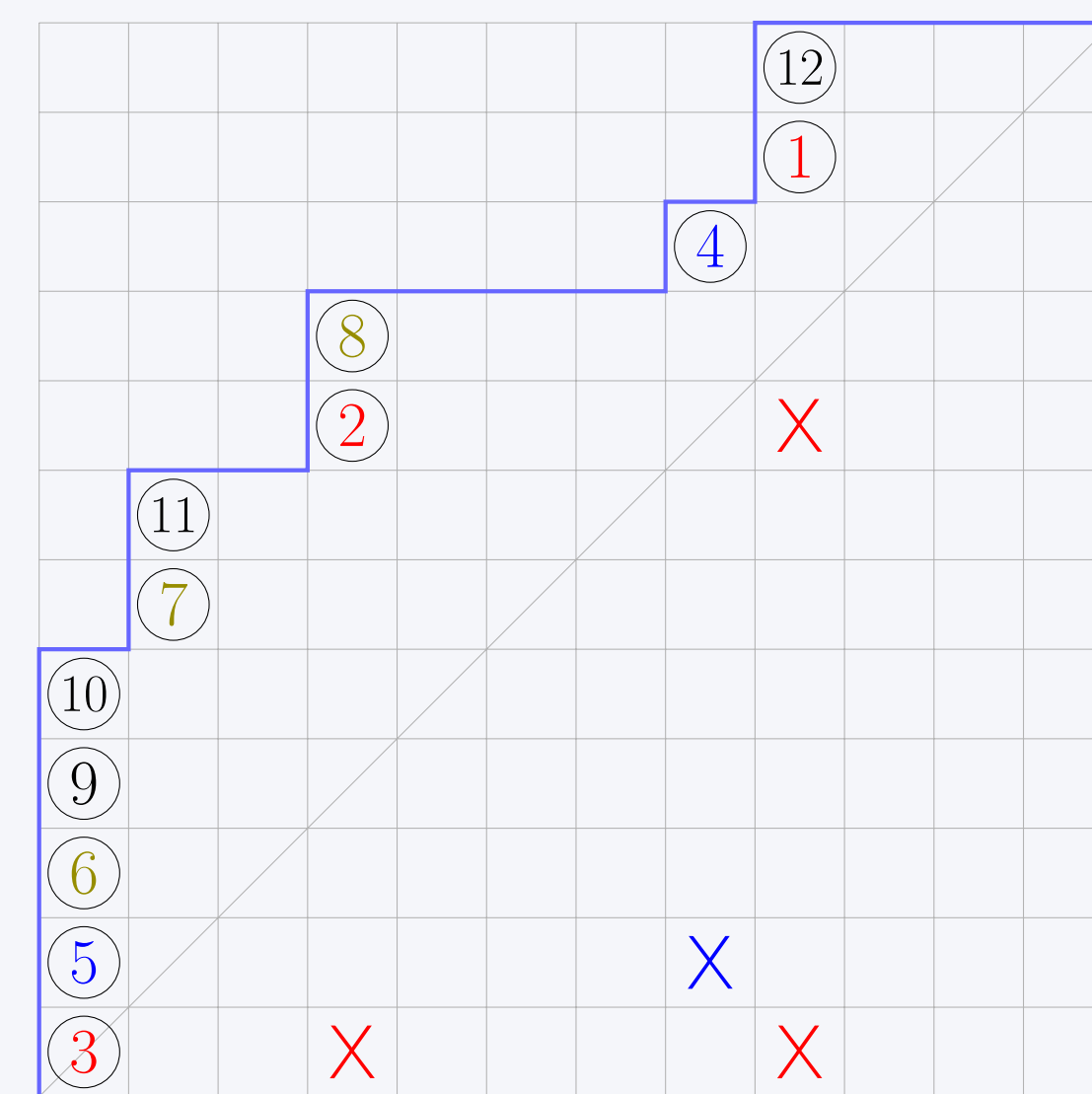
## Classical Statistic: level

$\text{level}(\kappa) = \text{total } \# \text{ grains} - \# \{ \text{non-sink edges} \}$ .

Previously known: level corresponds to area when  $\mu = (n)$  and  $\nu = \emptyset$ .

$\kappa(v_i) = j + \text{correction term} \iff j$  appears in  $i$ th column from the right.

$a_{12}(D) = 3$   
 $a_{11}(D) = 2$   
 $a_{10}(D) = 2$   
 $a_9(D) = 5$   
 $a_8(D) = 4$   
 $a_7(D) = 5$   
 $a_6(D) = 4$   
 $a_5(D) = 4$   
 $a_4(D) = 3$   
 $a_3(D) = 2$   
 $a_2(D) = 1$   
 $a_1(D) = 0$   
 $\text{area}(D) = 35$



## New Statistic: delay

$\text{delay}(\kappa)$  is computed via the following algorithm:

1. Topple 0
2. As long as some vertex hasn't been toppled, traverse the vertices of  $G$  in decreasing order.
3. When you arrive at an un-toppled vertex  $v$ :
  - if  $v$  can topple, topple it and continue.
  - If  $v$  cannot be toppled, increment the delay by 1 and continue.

## Computing delay

The configuration in our example has toppling word 10, 9, 7, 6, 5, 3, 2, 11, 8, 4, 1, 12, and hence a **delay** of 6. Following our algorithm step-by-step gives the following:

0. Fire 0.
1. In Round 1, we skip 12, 11, 8, 4, and 1.
2. In Round 2, we skip 12 again, but 11, 8, 4, and 1 topple.
3. In Round 3, 12 topples.

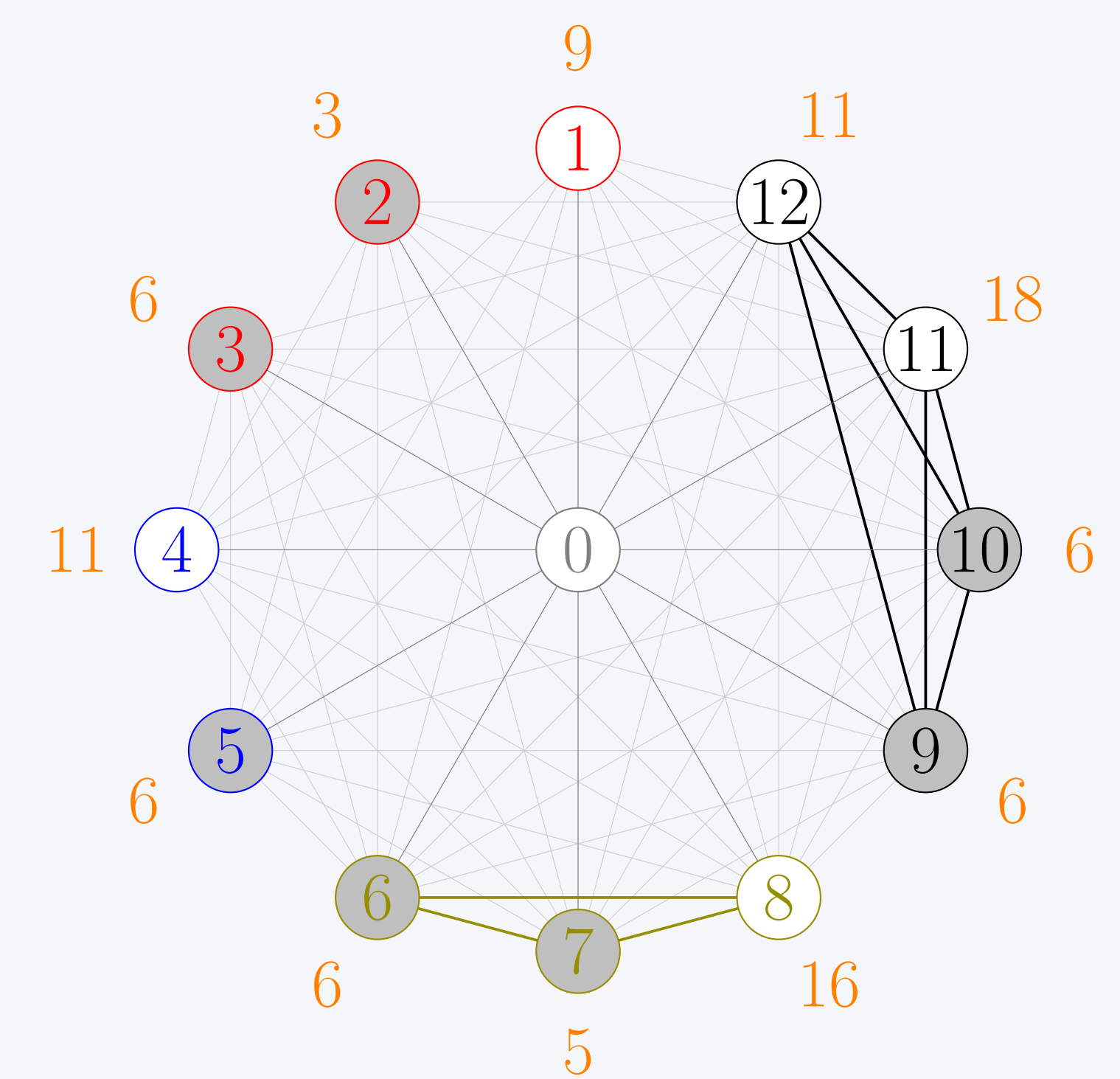


Figure 1. State of the configuration after Round 1. Nodes filled in gray have been toppled and are ignored in subsequent rounds.

## Main Result

The bivariate statistic  $(\text{level}, \text{delay})$  allows us to compute coeffs of  $\nabla e_n$  in terms of graphs.

**Theorem (DADILLBVW):**

For all  $n$  and  $\mu, \nu$  compositions with  $|\mu| + |\nu| = n$ , we have

$$\langle \nabla e_n, e_\mu h_\nu \rangle = \sum_{\kappa \in \text{sortrec}(\mu; \nu)} q^{\text{level}(\kappa)} t^{\text{delay}(\kappa)},$$

where  $\text{sortrec}(\mu; \nu)$  is the set of sorted recurrent configs on  $\hat{G}_{\mu, \nu}$ .

[1] Erik Carlsson and Anton Mellit, *A proof of the shuffle conjecture*, J. Amer. Math. Soc. **31** (2018), no. 3, 661–697. MR 3787405

[2] Nicholas A. Loehr and Jeffrey B. Remmel, *Conjectured combinatorial models for the Hilbert series of generalized diagonal harmonics modules*, Electron. J. Combin. **11** (2004), no. 1, Research Paper 68, 64. MR 2097334